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# Morphisms and Transformations of Potentially Inconsistent Graphs

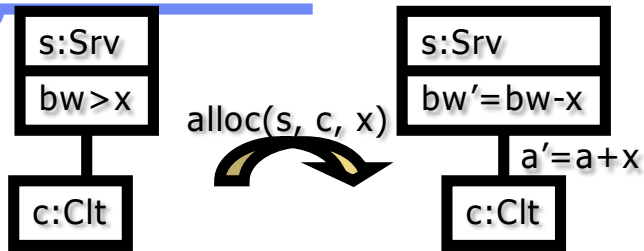
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Joint work with Fernando Orejas and Maryam  
Ghaffari Saadat

IFIP 1.3 Meeting 5-7 July 2018, Royal Holloway, UK



# Ground Attributed Graph Transformation



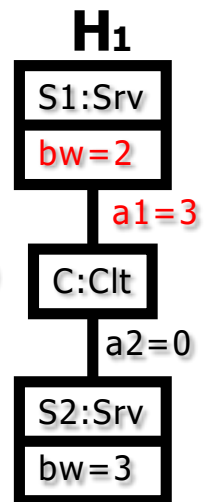
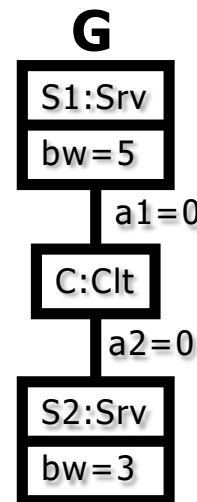
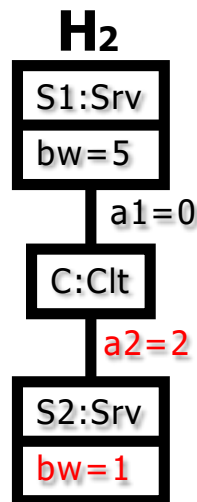
L, R attributed over  $\mathbf{X}$  with constraints

- $\Phi_L = \Phi_R = \{s.bw > x, s.bw' = s.bw - x, e.a' = e.a + x\}$
- Match  $m: L \rightarrow G$  satisfies  $D \models \Phi_G \rightarrow m(\Phi_L)$

→ Rule constraints satisfied by  $G$ 's attributes

Ground symbolic graphs  
SG = (G,  $\Phi$ )

- ✗ G attributed over vars  $\mathbf{X}$
- ✗ constraints  $\Phi = \{x = d, \dots\}$  with
  - vars  $x$  in  $\mathbf{X}$ ,
  - constants  $d$  in  $\Sigma$ -algebra  $\mathbf{D}$
- ✗ Invariant  $a1 + a2 < 4$

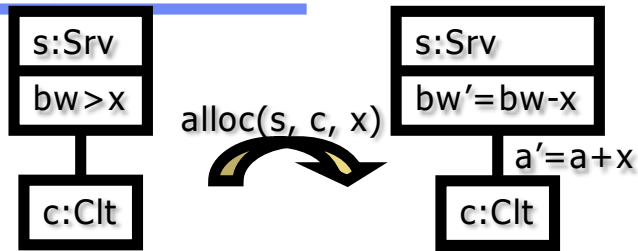


$\{S1.bw=5, a1=0, a2=2, S2.bw=1, a1+a2 < 4\}$

$\{S1.bw=5, a1=0, a2=0, S2.bw=3, a1+a2 < 4\}$

$\{S1.bw=2, a1=3, a2=0, S2.bw=3, a1+a2 < 4\}$

# Symbolic Attributed Graph Transformation



L, R attributed over  $\mathbf{X}$  with constraints

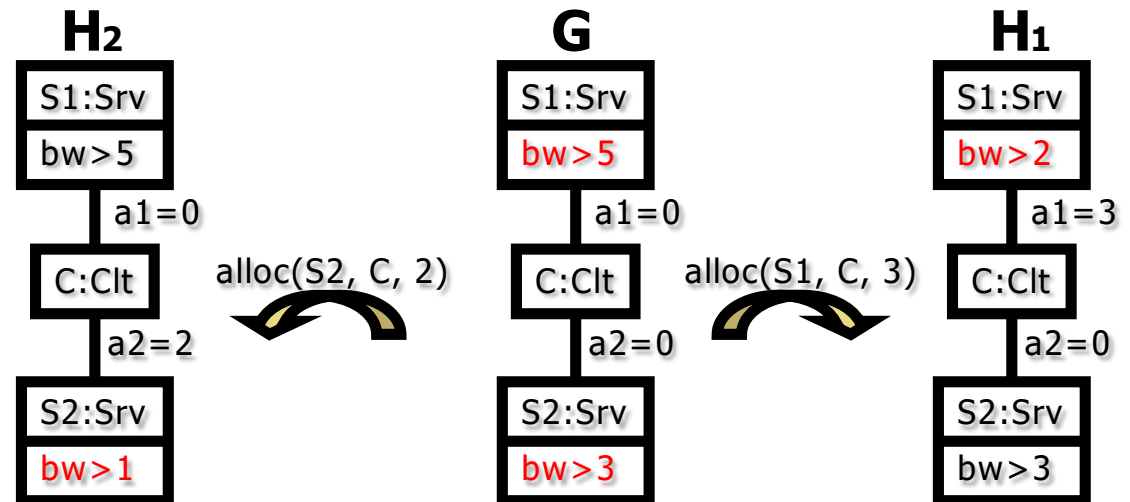
- $\Phi_L = \Phi_R = \{s.bw > x, s.bw' = s.bw - x, e.a' = e.a + x\}$
- Match  $m: L \rightarrow G$  satisfies  $D \models \Phi_G \rightarrow m(\Phi_L)$

→ Rule constraints entailed by G's constraints

(General) symbolic graphs  
 $SG = (G, \Phi)$

- ✗ G attributed over vars  $\mathbf{X}$
- ✗ FO constraints  $\Phi$  with
  - free vars in  $\mathbf{X}$  and
  - constants from  $\Sigma$ -algebra  $D$

$Sem(SG)$  = set of attr. graphs satisfying  $\Phi$



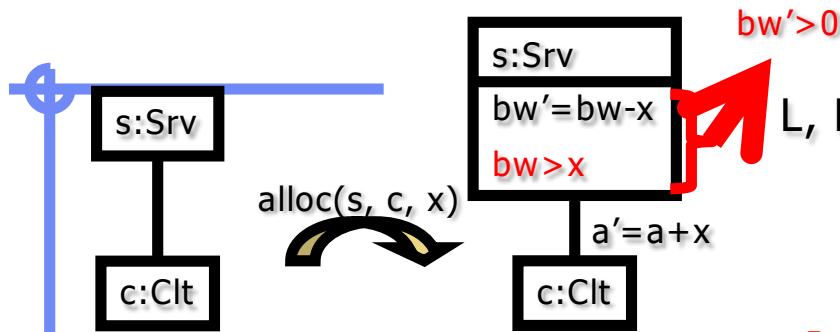
$\{S1.bw > 5, a1=0, a2=2, S2.bw > 1, a1+a2 < 4\}$

$\{S1.bw > 5, a1=0, a2=0, S2.bw > 3, a1+a2 < 4\}$

$\{S1.bw > 2, a1=3, a2=0, S2.bw > 3, a1+a2 < 4\}$

# Symbolic Attributed Graph Transformation

## with Narrowing



L, R attributed over  $X$  with constraints

- $\Phi_L = \{ \}$
- $\Phi_R = \{ bw' = bw - x, bw' > 0, e.a' = e.a + x \} \cup \Phi_L$
- Match  $m: L \rightarrow G$  satisfies  $D \models \Phi_G \rightarrow m(\Phi_L)$
- **Derived constraints  $\Phi_G \cup m_i(\Phi_R)$  consistent**

→ L's constraints entailed by G's constraints

→ R's constraints added to H, if consistent

(General) symbolic graphs

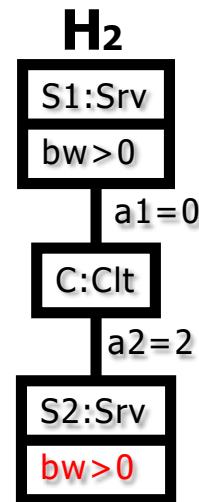
$SG = (G, \Phi)$

✗ G attributed over vars  $X$

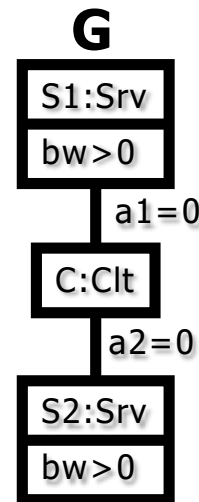
✗ FO constraints  $\Phi$  with

- free vars in  $X$  and
- constants from  $\Sigma$ -algebra  $D$

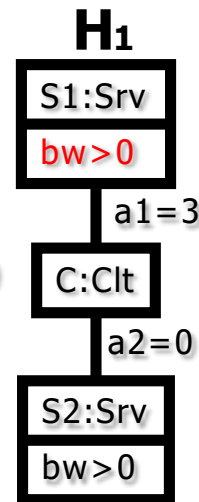
$Sem(SG)$  = set of attr. graphs satisfying  $\Phi$



$alloc(S2, C, 2)$



$alloc(S1, C, 3)$

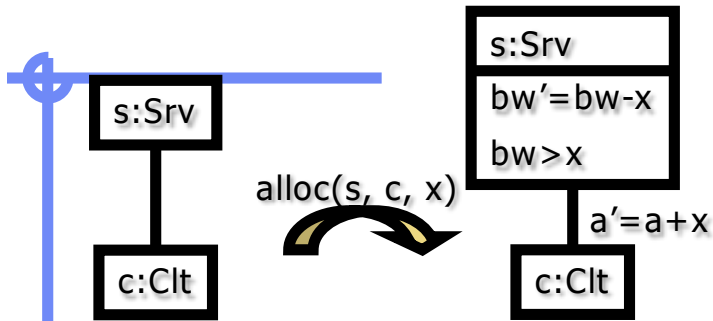


$\{S1.bw > 0, a1 = 0, a2 = 2, S2.bw > 0, a1 + a2 < 4\}$

$\{S1.bw > 0, a1 = 0, a2 = 0, S2.bw > 0, a1 + a2 < 4\}$

$\{S1.bw > 0, a1 = 3, a2 = 0, S2.bw > 0, a1 + a2 < 4\}$

# Symbolic Attributed Graph Transformation with Narrowing



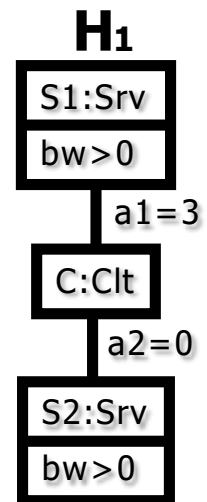
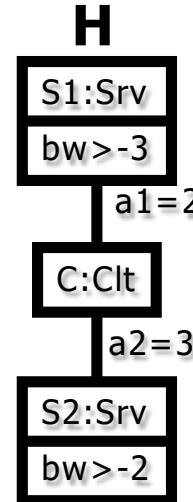
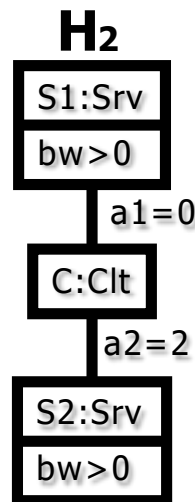
But: would like to

- ✗ decouple attribute handling from graph transformation, eg use external tool
- ✗ abstract representation of states with common graphs structure
- ✗ delay non-det. choices, retaining confluence

(General) symbolic graphs  
 $SG = (G, \Phi)$

- ✗  $G$  attributed over vars  $X$
- ✗ FO constraints  $\Phi$  with
  - free vars in  $X$  and
  - constants from  $\Sigma$ -algebra  $D$

$Sem(SG)$  = set of attr. graphs satisfying  $\Phi$



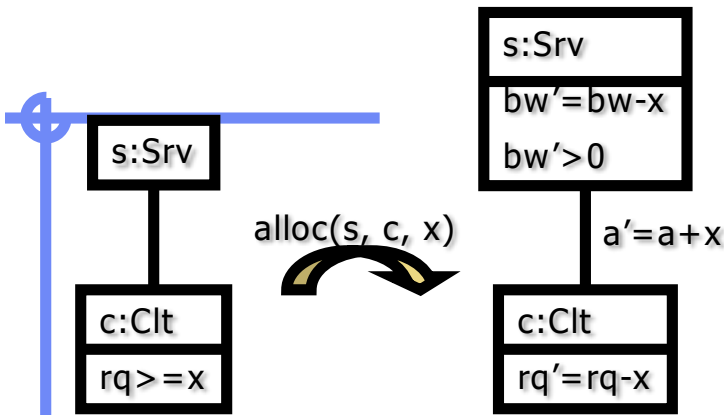
$\{S1.bw > 0, a1=0, a2=2, S2.bw > 0, a1+a2 < 4\}$

$\{S1.bw > -3, a1=2, a2=3, S2.bw > -2, a1+a2 < 4\}$

$\{S1.bw > 0, a1=3, a2=0, S2.bw > 0, a1+a2 < 4\}$

# Potentially Inconsistent Graph

## Transformation



L, R attributed over  $\mathbf{X}$  with constraints

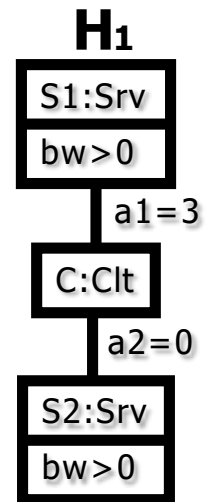
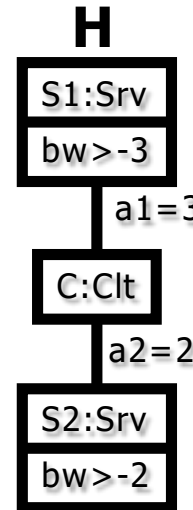
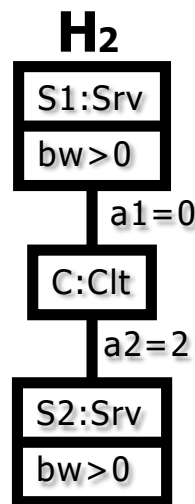
- $\Phi_L, \Phi_R$  as before
- Match  $m: L \rightarrow G$  satisfies  $D \models \Psi_G \rightarrow m(\Phi_L)$  for consistent subset  $\Psi_G$  of  $\Phi_G$

- ➔ L's constraints entailed by cons. weak. of G's
- ➔ R's constraints added to H, pot. inconsistent

PIGs = symbolic graphs  
 PIG =  $(G, \Phi)$

Pot(PIG) = all  $(G, \Psi)$  s.t.  $\Psi$   
 is max consist subset of  $\Phi$

Sem(PIG) =  $\bigcup_{(G, \Psi) \text{ in Pot(PIG)}} \text{Sem}(G, \Psi)$



$\{S1.bw > 0, a1=0, a2=2, S2.bw > 0, a1+a2 < 4\}$

$\{S1.bw > -3, a1=2, a2=3, S2.bw > -2, a1+a2 < 4\}$

$\{S1.bw > 0, a1=3, a2=0, S2.bw > 0, a1+a2 < 4\}$

# Any thoughts on ...

- ✗ Consistent notion of deduction?
  - $\Phi \vdash_c \varphi$  iff  $\Phi \vdash \varphi$  s.t. all formulas in the proof are consistent
  - See e.g. (Hunter & Nuseibeh 1998) on Inconsistency Management
- ✗ Institutions, categorical logic for para-consistency?
  - Is there a notion of morphism to get satisfaction condition, pushouts, etc?
- ✗ Precedents in state-based formal specifications, e.g., Z, B, ...?

# And now for the formal stuff...

- ✗ Morphisms
- ✗ Pushouts
- ✗ Institutions
- ✗ Local Church-Rosser